



Y. BEJARANO 3-16-1-49-1

SYSTEM AND METHOD FOR PROVISIONING QOS PATHS WITH RESTORATION IN A NETWORK

ATTORNEY: DAVID H. HITT (972) 480-8800

REPLACEMENT SHEET

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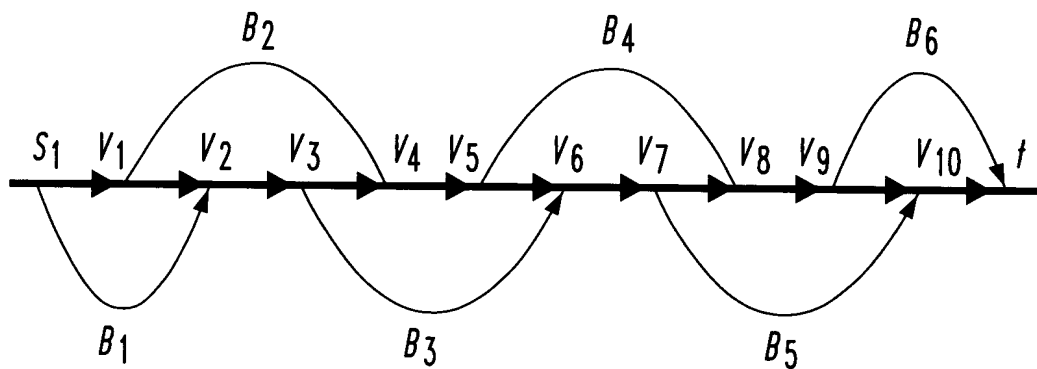


FIG. 1

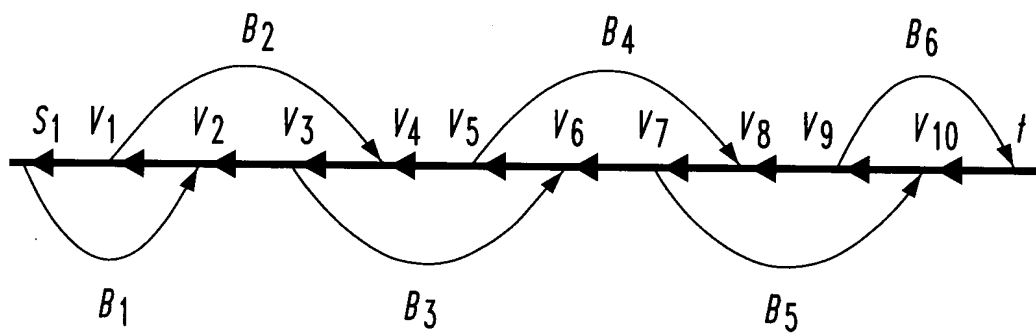


FIG. 2

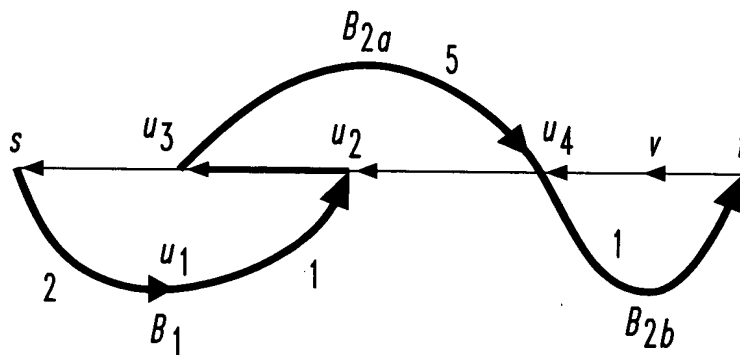


FIG. 3



FIG. 4

Algorithm PP (  $G(V, E), \hat{P}, \hat{d}, U$  )

parameters:

$G(V, E)$  - network,

$\{d_t, c_t\}_{t \in E}$  - delays and costs of the network links,

$\hat{P} = \{s = v_0, v_1, \dots, t = v_n\}$  - QoS path,

$\hat{d}$  - delay constraint,

$U$  - the upper bound on the cost of  $\mathcal{R}$ .

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1   $\Delta \leftarrow \hat{d} - D(\hat{P})$ 
2   $E' \leftarrow E$ 
3  for each link  $l = (v_i, v_{i+1}) \in \hat{P}$  do
4     $E' \leftarrow E' \setminus \{(v_i, v_{i+1}) \in \hat{P}\}$ 
5     $E' \leftarrow E' \cup \{(v_{i+1}, v_i) \in \hat{P}\}, c_{(v_{i+1}, v_i)} \leftarrow 0$ 
6  for all  $v_i \in V$  do
7     $D_{v_i}[0] \leftarrow \infty$ 
8     $D_s[0] \leftarrow 0$ 
9  for  $c = 1, 2, \dots, U$  do
10   for each  $v_j \in V$  in order such that  $v_j$  is before  $v_{j'}$  if  $v_j$  is a successor of  $v_{j'}$  in  $\hat{P}$  do
11      $D_{v_j}[c] \leftarrow D_{v_j}[c-1]$ 
12   for each link  $l = (v_i, v_j) \in E'$  do
13     RELAX( $l(v_i, v_j), c, \Delta$ )
14   if  $D_t[c] \leq D(\hat{P})$  then
15     determine walk  $\mathcal{W}$  by backtracking
16     return  $\mathcal{W}$ .
17 return FAIL

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FIG. 4

— CONTINUED FROM SHT 2/8

Procedure RELAX ( $l = (v_i, v_j), c, \Delta$ )

- 1 if  $v_j \in \hat{\mathcal{P}}$  and  $v_i \in \hat{\mathcal{P}}$  then
- 2 if  $D_{v_i}[c] \leq D(\hat{\mathcal{P}}_{(s, v_i)})$  then
- 3  $D_{v_j}[c] \leftarrow \min\{D_{v_j}[c], D(\hat{\mathcal{P}}_{(s, v_j)})\}$
- 4 else
- 5 if  $c_l \leq c$  then
- 6  $D_{v_j}[c] \leftarrow \min\{D_{v_j}[c], D_{v_i}[c - c_l] + d_l\}$
- 7 if  $v_j \in \hat{\mathcal{P}}$  and  $D_{v_j}[c] \leq D(\hat{\mathcal{P}}_{(s, v_j)}) + \Delta$  then
- 8  $D_{v_j}[c] \leftarrow \min\{D_{v_j}[c], D(\hat{\mathcal{P}}_{(s, v_j)})\}$

FIG. 5

Algorithm RT (  $G(V, E), \hat{P}, \hat{d}, \varepsilon$  )

parameters:

$G(V, E)$  - network

$\hat{P} = \{s = v_1, v_2, \dots, t = v_n\}$  - QoS path,

$\hat{d}$ - delay constraint

$\varepsilon$ - approximation ratio

```

1   $L, U \leftarrow \text{BOUND}(G(V, E), \hat{P}, \hat{d})$ 
2  do
3     $B \leftarrow \sqrt{L \cdot U}$ 
4    if TEST( $G(V, E), \hat{P}, \hat{d}, B, \varepsilon$ ) returns YES then
5       $L \leftarrow B$ 
6    else
7       $U \leftarrow 2 \cdot B$ 
8    until  $U/L \leq 8$ .
9   $\mathcal{W} \leftarrow \text{SCALE}(G(V, E), \hat{P}, \hat{d}, L, U, \varepsilon)$ 
10 return the restoration topology that corresponds to  $\mathcal{W}$ .

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Procedure SCALE( $G(V, E), \hat{P}, \hat{d}, L, U, \varepsilon$ )

```

1   $S \leftarrow \frac{L\varepsilon}{2N}$ 
2  for each link  $l \in E$  do
3     $c'_l \leftarrow \lfloor \frac{c_l}{S} \rfloor + 1$ 
4     $\tilde{U} \leftarrow \lfloor \frac{U}{S} \rfloor + 2N$ 
5  return PP( $G(V, E), \{d_i, c'_i\}_{i \in E}, \hat{P}, \hat{d}, \tilde{U}$ )

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FIG. 5

— CONTINUED FROM SHT 4/8

Procedure TEST( $G(V, E), \hat{P}, \hat{d}, B$ )

- 1 Apply Procedure SCALE for ( $G(V, E), \hat{P}, \hat{d}, B, 2$ )
- 2 if Algorithm SCALE returned FAIL then
- 3     return NO
- 4 else
- 5     return YES

Procedure BOUND( $G(V, E), \hat{P}, \hat{d}$ )

- 1 let  $c^1 < c^2 < \dots < c^r$  the distinct costs values of the links.
- 2  $low \leftarrow 1$ ;  $high \leftarrow r$
- 3 while  $low < high - 1$
- 4      $j \leftarrow \lfloor (high + low) / 2 \rfloor$
- 5      $E' \leftarrow \{l \mid c_l \leq c^j\}$
- 6     set  $c_l \leftarrow 1$  for each  $l \in E'$
- 7     apply Algorithm PP on ( $G'(V, E'), \hat{P}, \hat{d}, 2N$ )
- 8     if Algorithm PP returned FAIL then
- 9          $high \leftarrow j$
- 10    else
- 11        $low \leftarrow j$
- 12  $U \leftarrow 2N \cdot c^{high}$ ;  $L \leftarrow c^{high}$ ;
- 13 return  $L, U$ ;

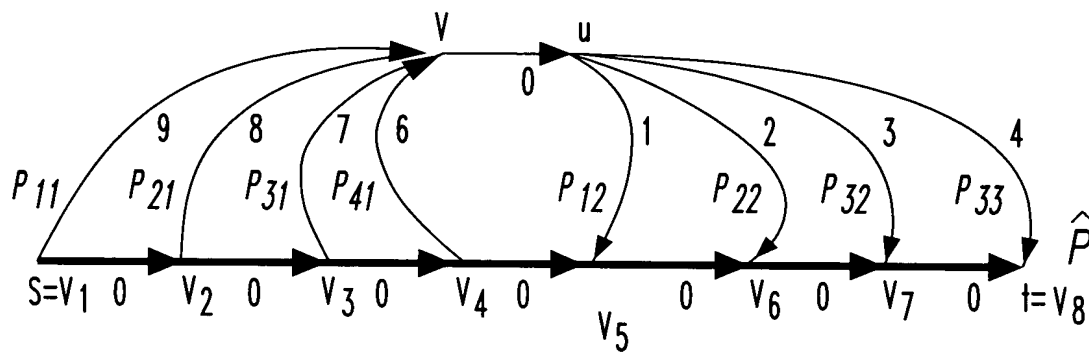
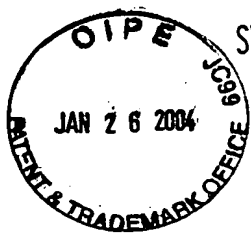


FIG. 6

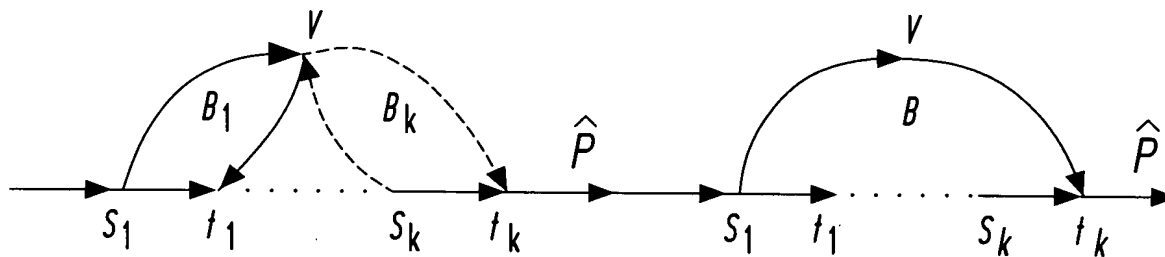
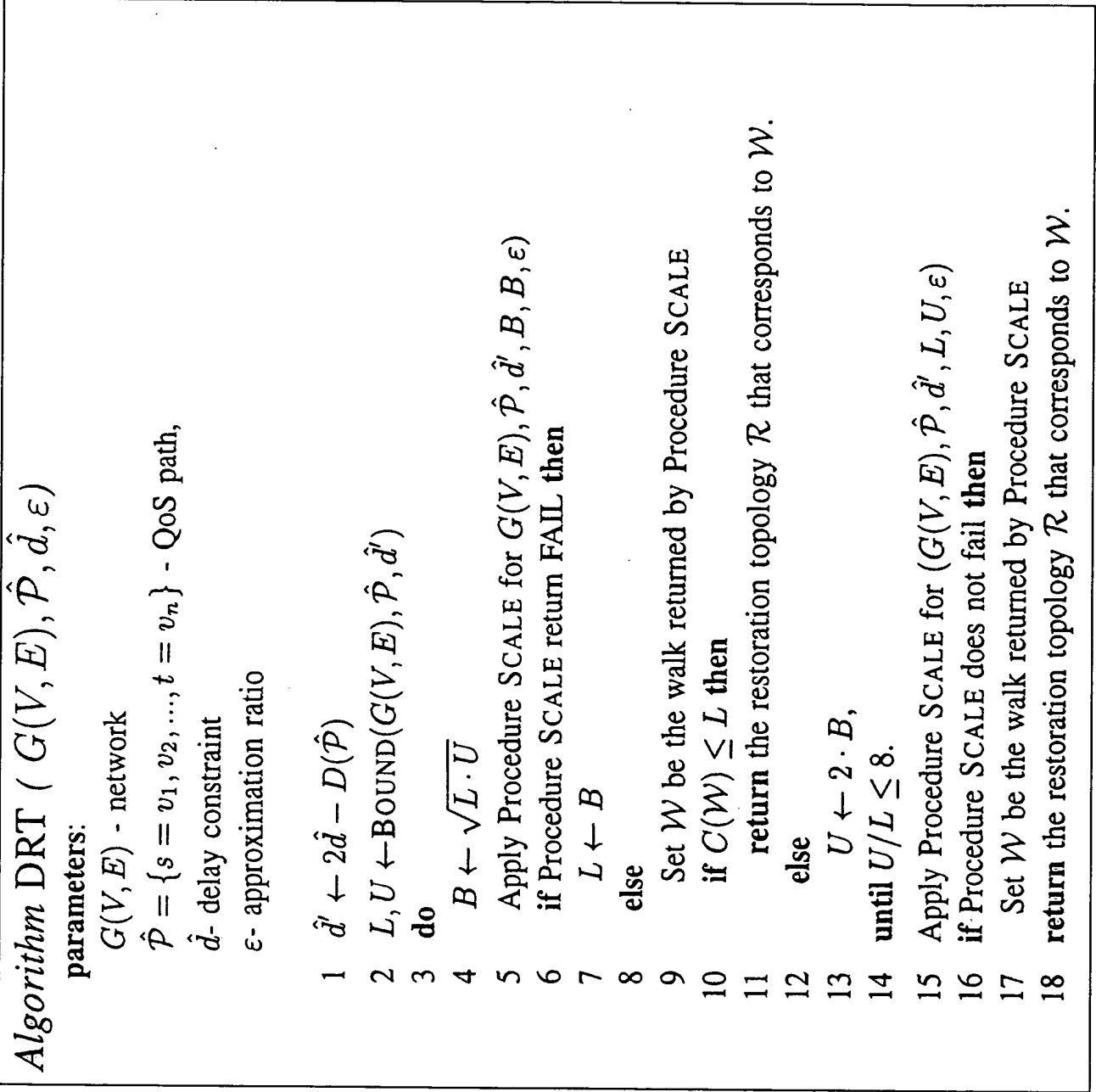


FIG. 7



FIG. 8



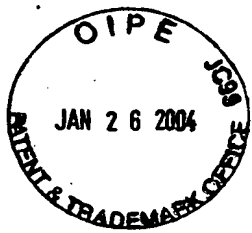


FIG. 9

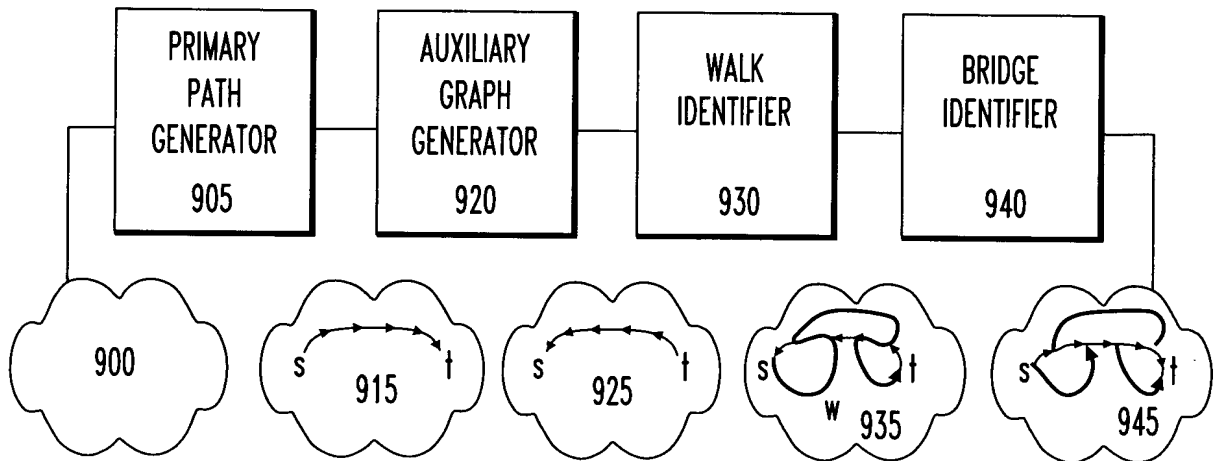


FIG. 10

